Quantified Boolean Formulas and DepQBF

4. DepQBF in Practice

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Overview

- DepQBF: search-based, QCDCL solver.
- First release in February 2010, under active development.
- Approx. 20,000 lines of C code.
- Open source under GPL: http://lonsing.github.io/depqbf/
- "DepQBF": optional dependency analysis to relax the quantifier ordering.
- Design decision: allow for use as a library.
- No pre/inprocessing (not yet officially released).
- Trace generation for certificate generation.
- Based on PCNF, QDIMACS input format.
- Incremental solving: beneficial when solving sequences of closely related PCNFs.
- API to manipulate the input PCNF, configure the solver.
- New version 4.0 released in February 2015.

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Input Format

QDIMACS:

- Extension of DIMACS format used in SAT solving.
- Easy to parse.
- Literals of variables encoded as signed integers.
- One quantifier block per line ("a" labels \forall , "e" labels \exists), terminated by zero.
- One clause per line, terminated by zero.

Example $\exists x_1, x_3, x_4 \forall y_5 \exists x_2.(\neg x_1 \lor x_2) \land (x_3 \lor y_5 \lor \neg x_2) \land (x_4 \lor \neg y_5 \lor \neg x_2) \land (\neg x_3 \lor \neg x_4) = \begin{cases} p \text{ cnf } 5 \text{ 4} \\ e \text{ 1 } 3 \text{ 4 } 0 \\ a \text{ 5 } 0 \\ e \text{ 2 } 0 \\ e \text{ 2 } 0 \\ -1 \text{ 2 } 0 \\ 3 \text{ 5 } -2 \text{ 0} \\ 4 \text{ -5 } -2 \text{ 0} \\ -3 \text{ -4 } 0 \end{cases}$

Using DepQBF in Your Application

- Encode your problem in QDIMACS format: support for other formats?
- DepQBF is a standalone QBF solver and...
- ... provides a library with a API in C: add a formula, solve, ...
- Library use is more convenient: incremental calls.

- Compile DepQBF, which produces the library libqdpll.a.
- Include the header file qdpll.h in your source code.
- Compile and link against the solver library: gcc your_code.c -L. -lqdpll
- Call the solver API from your application.

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```
/* Create and initialize solver instance. */
QDPLL * qdpll_create (void);
```

/* Delete solver instance and release all memory. */
void qdpll_delete (QDPLL * qdpll);

/* Ensure variable table size to be at least 'num'. */
void qdpll_adjust_vars (QDPLL * qdpll, VarID num);

- No static data: allows to generate multiple solver objects.
- DepQBF uses variable indices as given by the QDIMACS file to index a table of variable objects: keep indices compact in the encoding.

/* Configure solver instance via configuration string. Returns null pointer on success and error string otherwise. */ char * qdpll_configure (QDPLL * qdpll, char * configure_str);

Possible configuration strings:

- Call ./depqbf -h for a partial listing of options.
- --no-cdcl: disable clause learning and backtrack chronologically from conflicts.
- --no-sdc1: disable cube learning backtrack chronologically from solutions.
- --no-pure-literals: disable pure literal detection.
- Various learning variants: long-distance resolution, lazy learning.
- Many more: heuristics,...

Prefix Manipulation:

- Add quantifier blocks of any type at any prefix position.
- Add new variables to quantifier blocks.
- No explicit deletion of blocks/variables: garbage collection.

CNF Manipulation:

- Add/delete clauses.
- No modifications of present clauses: must delete old and add new clause.

Stack-Based Clause Additions/Deletions:

- Push new clauses onto the clause stack.
- Pop most recently added clauses from the stack.

- Declare groups (sets) of clauses.
- Add clauses to a group, remove entire groups.

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API: Prefix Manipulation (1/3)

```
enum QDPLLQuantifierType:
  QDPLL_QTYPE_EXISTS = -1
  QDPLL_QTYPE_UNDEF = 0
  QDPLL_QTYPE_FORALL = 1
```

typedef unsigned int Nesting;

/* Add new quantifier block with type 'qtype' at right end of prefix. */
Nesting qdpll_new_scope (QDPLL * qdpll, QDPLLQuantifierType qtype);

API: Prefix Manipulation (2/3)

typedef unsigned int VarID;

```
/* Add new variable 'id' to the block at level 'nesting'.
Fails if a variable with 'id' already exists. */
void qdpll_add_var_to_scope (QDPLL * qdpll, VarID id, Nesting nesting);
```

typedef int LitID;

/* Add new variable 'id' to the current quantifier block
 opened by a previous call of 'qdpll_new_scope' or
 'qdpll_new_scope_at_nesting'.
 Adding '0' closes the current block.
 Fails if a variable with 'id' already exists. */
void qdpll_add (QDPLL * qdpll, LitID id);

/* Returns the nesting level of the current rightmost block. */
Nesting qdpll_get_max_scope_nesting (QDPLL * qdpll);

/* Return largest declared variable ID. */
VarID qdpll_get_max_declared_var_id (QDPLL * qdpll);

/* Returns non-zero iff. variable 'id' has been added to the formula. */
int qdpll_is_var_declared (QDPLL * qdpll, VarID id);

/* Return nesting of block which contains variable 'id'. */
Nesting qdpll_get_nesting_of_var (QDPLL * qdpll, VarID id);

/* Return the type of the block at level 'nesting'.*/
QDPLLQuantifierType qdpll_get_scope_type (QDPLL *qdpll, Nesting nesting);

API: Basic CNF Manipulation

/* Add a literal 'id' to the current open clause. Adding '0' closes the clause. */ void qdpll_add (QDPLL * qdpll, LitID id);

/* Pretty-print PCNF to 'out' using QDIMACS format. */
void qdpll_print (QDPLL * qdpll, FILE * out);

- Note: qdpll_add is used to add variables to blocks and literals to clauses.
- Tautological input clauses are discarded.
- Superfluous literals (double occurrences) in clauses are discarded.
- Literals in input clauses are sorted by prefix order and universal-reduced.
- No free variables: if id in a clause is a literal of new variable, then that variable is put into a default existential quantifier block $\exists B_0$ at the left end of the prefix: $\exists B_0 Q_1 B_1 \dots Q_n B_n . \phi$.
- In practice: first add the prefix, then the clauses.

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- In practice: first add the prefix, then the clauses.

/* Open a new top-most frame on the clause stack. Clauses added by 'qdpll_add' are added to the top-most frame. */ unsigned int qdpll_push (QDPLL * qdpll);

/* Pop the top-most frame from the clause stack. The clauses in that frame are considered deleted from the formula. */ unsigned int qdpll_pop (QDPLL * qdpll);

/* Enforce garbage collection of popped off clauses. */
void qdpll_gc (QDPLL * qdpll);

API: Stack-Based CNF Manipulation (2/2)

- Must configure by --dep-man=simple (use given linear quantifier ordering)...
- ...and by --incremental-use.
- Useful if a sequence of closely related PCNFs is solved.
- Example: encoding a transition relation for i steps, i + 1 steps,...
- No need to parse all the PCNFs from scratch, but only the new clauses.
- More important: learned clauses and cubes are re-used (if possible) when solving other PCNFs in the sequence.
- Incorrect learned clauses and cubes are automatically discarded.
- Pushing is optional: without any push before, clauses are added to a default frame and cannot be removed.

In Practice:

- Push and add clauses which are shared between the PCNFs first.
- Push clauses which have to be removed last, so that they can be deleted by a pop.

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API: Group-Based CNF Manipulation (1/3)

/* Declare a new clause group and return its handle. */
ClauseGroupID qdpll_new_clause_group (QDPLL *qdpll);

/* Delete the clause group 'cg'. Its handle 'cg' becomes invalid. */
void qdpll_delete_clause_group (QDPLL *qdpll, ClauseGroupID cg);

/* Returns non-zero iff. a clause group with handle 'cg' exists. */
int qdpll_exists_clause_group (QDPLL *qdpll, ClauseGroupID cg);

/* Open the clause group 'cg'. Clause can be added to open groups only. Only one group can be open at a time. */ void qdpll_open_clause_group (QDPLL *qdpll, ClauseGroupID cg);

/* Closes the open clause group 'cg'. */
void qdpll_close_clause_group (QDPLL *qdpll, ClauseGroupID cg);

/* Returns handle of currently open clause group or zero if there is no open group. */ ClauseGroupID qdpll_get_open_clause_group (QDPLL *qdpll);

API: Group-Based CNF Manipulation (2/3)

/* Temporarily disable (remove) clauses in clause group 'cg'. */
void qdpll_deactivate_clause_group (QDPLL *qdpll, ClauseGroupID cg);

/* If the formula is unsatisfiable (qdpll_sat() == QDPLL_RESULT_UNSAT):
 returns zero-terminated list of clause groups containing clauses
 which were actually used by the solver (unsatisfiable core). */
ClauseGroupID * qdpll_get_relevant_clause_groups (QDPLL * qdpll);

/* Enforce garbage collection of deleted groups. */
void qdpll_gc (QDPLL * qdpll);

API: Group-Based CNF Manipulation (3/3)

- Must configure by --dep-man=simple (use given linear quantifier ordering)...
- ...and by --incremental-use.
- More flexible than stack-based API (push/pop).
- Calls of stack/group-based API functions cannot be mixed.
- Allows deletion of arbitrary clauses instead of only most recently added ones.

In Practice:

- In general, clause groups result in increased overhead compared to push/pop API.
- Use push/pop whenever appropriate.

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API: Deletion of Clauses, Variables, and Quantifier Blocks

Clauses:

- No explicit deletion through API.
- A clause is considered deleted after its frame has been popped from the stack.
- A clause is considered deleted after its group has been deleted.
- Garbage collection triggered heuristically, or enforced by calling qdpl1_gc.

Variables:

- No explicit deletion through API.
- A variable x is deleted by qdpll_gc if all the clauses where x occurs have been deleted before.
- The IDs of deleted variables can be re-used: check with qdpll_is_var_declared.

Quantifier Blocks:

- No explicit deletion through API.
- A quantifier block is deleted by qdpll_gc if all of its variables have been deleted before.

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API: Solving (1/2)

```
enum QDPLLResult:
   QDPLL_RESULT_UNKNOWN = 0
   QDPLL_RESULT_SAT = 10
   QDPLL_RESULT_UNSAT = 20
```

/* Solve the given PCNF. */
QDPLLResult qdpll_sat (QDPLL * qdpll);

/* Reset internal solver state, but keep the PCNF and learned constraints.
void qdpll_reset (QDPLL * qdpll);

```
/* Discard all learned constraints. */
void qdpll_reset_learned_constraints (QDPLL * qdpll);
```

- QDPLL_RESULT_UNKNOWN returned only if formula not solved under imposed limits.
- qdpll_reset deletes the variable assignments.
- Incremental calls after reset: may add/delete clauses.
- For convenience: calling qdpll_reset_learned_constraints is never required for the correctness of incremental solving. The solver keeps track of learned constraints.

API: Solving (2/2)

typedef int QDPLLAssignment;
#define QDPLL_ASSIGNMENT_FALSE -1
#define QDPLL_ASSIGNMENT_UNDEF 0
#define QDPLL_ASSIGNMENT_TRUE 1

/* Get current assignment of variable. */
QDPLLAssignment qdpll_get_value (QDPLL * qdpll, VarID id);

/* Like 'qdpll_get_value' but print to standard output. */
void qdpll_print_qdimacs_output (QDPLL * qdpll);

- Call after qdpll_sat but before qdpll_reset.
- From the command line: --qdo
- Get partial certificates of (un)satisfiability as assignments to leftmost variables...
- ... if the PCNF $\exists B_1 \ldots, \phi$ is satisfiable.
- ... if the PCNF $\forall B_1 \dots, \phi$ is unsatisfiable.
- In practice: useful for encodings of problems from the second level of the polynomial hierarchy with prefix ∀∃ and ∃∀.

API: Solving Under Assumptions

/* Assign a variable permanently in the next run (assumption).
 If 'id < 0' then assign variable 'id' to false.
 If 'id > 0' then assign variable 'id' to true. */
void qdpll_assume (QDPLL * qdpll, LitID id);

/* Returns an array of safe arguments to 'qdpll_assume'. */
LitID * qdpll_get_assumption_candidates (QDPLL * qdpll);

/* Returns the subset of assumptions used by the solver to determine the result. */

LitID * qdpll_get_relevant_assumptions (QDPLL * qdpll);

• Can be combined with push and pop.

- Safe arguments to qdpll_assume are variables from the leftmost block (recursively).
- Assignments added by qdpll_assume are persistent in the next call of qdpll_sat.
- qdpll_reset removes assignments added by qdpll_assume before.
- Constraints learned under assumptions are correct independently.
- For convenience: calling qdpll_reset_learned_constraints is never required for the correctness of incremental solving. The solver keeps track of learned constraints.

API: Generating Traces and Certificates

/* Configure solver instance via configuration string. Returns null pointer on success and error string otherwise. */ char * qdpll_configure (QDPLL * qdpll, char * configure_str);

- Print the *full* resolution derivation in QRP format to standard output: can be huge!
- --trace=qrp (text format) or --trace=bqrp (binary format).
- QBFcert framework: http://fmv.jku.at/qbfcert/.
- Acknowledgments: Aina Niemetz and Mathias Preiner.
- Resolution proof checking by QRPcheck: http://fmv.jku.at/qrpcheck/.
- Certificate extraction (Skolem/Herbrand functions) by QRPcert: http://fmv.jku.at/qrpcert/.
- Skolemization/Herbrandization by CertCheck: http://fmv.jku.at/certcheck/.
- Checking skolemized/herbrandized formula using a SAT solver.

Remarks

- Please publish your benchmarks!
- Effective use of QBF solvers (sometimes) requires expert knowledge and fine tuning.
- Long-term goal: usability, integrated workflow.